# Sparse Representation and Dictionary Learning



Suh-Yuh Yang (楊肅煜)

Department of Mathematics, National Central University Jhongli District, Taoyuan City 320317, Taiwan

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#### **Outlines**

- Sparse representation and dictionary learning
  - Sparse representation (SR) problem
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# Part I

**Sparse Representation and Dictionary Learning** 

# Sparse representation problem

Terms: Sparse Representation (稀疏表現)/Sparse Coding (稀疏編碼)

**SR problem:** Given a signal vector  $\mathbf{x} \in \mathbb{R}^m$  and a dictionary matrix  $\mathbf{D} \in \mathbb{R}^{m \times n}$ , we seek a sparse coefficient vector  $\mathbf{z}^* \in \mathbb{R}^n$  such that

$$z^* = \arg\min_{z} \left(\frac{1}{2} \|x - Dz\|_2^2 + \lambda \|z\|_0\right),$$

where  $\lambda > 0$  is a penalty parameter and  $\|z\|_0$  counts the number of nonzero components of z.

#### Remarks:

- In the matrix-vector multiplication *Dz*, the components of *z* are the coefficients with respect to columns (also called *atoms*) of *D*.
- We call  $||z||_0$  the  $\ell^0$  norm of z, even though  $\ell^0$  is *not* really a norm, since the *homogeneity property* fails,  $||\alpha z||_0 \neq |\alpha|||z||_0$ .
- It is inefficient to compute  $||z||_0$  directly when n is large. In practice, we will use the  $\ell^1$  norm instead of the  $\ell^0$  norm.

# Two dual $\ell^0$ minimization problems

In [Sharon-Wright-Ma 2007], they studied the following two dual  $\ell^0$  minimization problems:

• **Sparse error correction (SEC):** Given  $0 \neq y \in \mathbb{R}^n$  and  $A \in \mathbb{R}^{n \times p}$  with n > p and rank(A) = p, we seek  $w^* \in \mathbb{R}^p$  such that

$$w^* = \underset{w}{\operatorname{arg\,min}} \| \boldsymbol{y} - \boldsymbol{A} \boldsymbol{w} \|_0. \tag{1}$$

• Sparse signal reconstruction (SSR): Given  $D \in \mathbb{R}^{m \times n}$  with m < n and  $0 \neq x \in C(D)$  the column space of D, we seek  $z^* \in \mathbb{R}^n$  such that

$$z^* = \underset{z}{\operatorname{arg\,min}} \|z\|_0$$
 subject to  $x = Dz$ . (2)

Note that (1) is a *decoding problem*, while (2) is a *sparse representation problem*. These two problems are *dual* in the sense that we can convert one problem to the other, see page 8 below.

## Existence and uniqueness of solution

- Existence:
  - Existence of  $w^*$ : If  $\exists w \in \mathbb{R}^p$  s.t.  $\|y Aw\|_0 = 0$ , then  $w^* = w$ . Otherwise, define

$$\mathcal{S} := \{ k \in \mathbb{N} : \exists \ \mathbf{w} \in \mathbb{R}^p \text{ s.t. } \|\mathbf{y} - \mathbf{A}\mathbf{w}\|_0 = k \}.$$

Then  $\emptyset \neq S \subseteq \mathbb{N}$ . By the well-ordering principle,  $\exists k_0 \in S$  the minimum of S. i.e.,  $\exists w^*$  such that  $w^* = \arg\min_{z_0} ||y - Aw||_0$ .

- Existence of  $z^*$ : It can be shown in a similar way!
- **2 Uniqueness:** *It will generally be true* that these two dual problems have a unique solution if
  - $\exists w_0$  such that the error  $e := y Aw_0$  is sparse enough, or
  - $\exists$   $z_0$  sparse enough such that  $x = Dz_0$ . e.g., if any set of 2T columns of D are linearly independent, then any  $z_0 \in \mathbb{R}^n$  with  $||z_0||_0 \le T$  such that  $Dz_0 = x$  is the unique solution to SSR problem (2).

## Why we require matrix A full rank p in the SEC problem?

Note that *A* is of size  $n \times p$  and n > p.

Suppose that *A* is not full rank *p*. Then rank(A) < p.

Since  $\dim N(A) + \operatorname{rank}(A) = p$ , we have  $\dim N(A) > 0$ .

Thus, nullspace  $N(A) \neq \{\mathbf{0}\}$  and  $\exists \ \widetilde{w} \neq \mathbf{0}$  such that  $A\widetilde{w} = \mathbf{0}$ .

If  $w^*$  is a solution of the SEC problem, then

$$\|y - A(w^* + \widetilde{w})\|_0 = \|y - Aw^*\|_0.$$

Hence,  $w^* + \widetilde{w}$  is also a solution of the SEC problem.

Therefore, in order to ensure the uniqueness, we require A full rank p.

### How to convert problem (2) to problem (1)?

- The decoding problem (1) can be converted to the sparse representation problem (2). [Candès et al. 2005, IEEE Symposium on FOCS]
- Converting (2) to (1): Let p = n rank(D) > 0 and A be a full-rank  $n \times p$  matrix whose columns span the nullspace of D, i.e., DA = 0. Find any  $y \in \mathbb{R}^n$  so that Dy = x and define f(w) = y Aw. Then

$$\underset{z^*}{\arg\min} \frac{\|z\|_0}{\underbrace{\sum \underbrace{\& Dz = x}}} = \underbrace{f\left(\arg\min_{w} \|y - Aw\|_0\right)}_{f(w^*)}. \tag{3}$$

*Proof:* First, note that for all  $w \in \mathbb{R}^p$ , we have

$$Df(w) = D(y - Aw) = Dy - DAw = Dy = x.$$

Claim:  $\exists \ \widetilde{w} \in \mathbb{R}^p$  such that  $f(\widetilde{w}) = y - A\widetilde{w} = z^*$ .

$$\therefore Dz^* = x \text{ and } D(y - Aw) = x, \ \forall w \implies D(-z^* + y - Aw) = 0$$

$$\therefore$$
  $\exists$   $\bar{w}$  such that  $A\bar{w}=-z^*+y-Aw \implies z^*=y-A(w+\bar{w}):=f(\tilde{w})$ 

Claim:  $\widetilde{w} = w^* := \arg\min_{w} \|y - Aw\|_0$ , and then  $f(w^*) = z^*$ .

$$\ \, :: \|f(\boldsymbol{w}^*)\|_0 \leq \|f(\widetilde{\boldsymbol{w}})\|_0 = \|z^*\|_0 \leq \|f(\boldsymbol{w}^*)\|_0 \Longrightarrow \|f(\boldsymbol{w}^*)\|_0 = \|f(\widetilde{\boldsymbol{w}})\|_0$$

By the uniqueness of  $w^*$ , we obtain  $\widetilde{w} = w^*$  and then  $f(w^*) = z^*$ .

# The $\ell^1$ - $\ell^0$ equivalence problem

• In general, the  $\ell^0$  minimizations (1) and (2) are *NP-hard problems*:

$$w^* = \underset{w}{\operatorname{arg\,min}} \| y - Aw \|_0, \tag{1}$$

$$z^* = \underset{z}{\arg\min} \|z\|_0$$
 subject to  $x = Dz$ . (2)

• The equivalence between  $\ell^0$  and  $\ell^1$  minimizations is conditional.

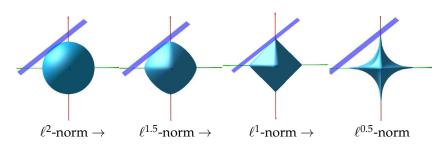
David L. Donoho, For most large underdetermined systems of linear equations the minimal  $\ell_1$ -norm solution is also the sparsest solution, CPAM, 59 (2006), pp. 797-829.

If the error  $e := y - Aw^*$  or the solution  $z^*$  is sufficiently sparse, then the solutions to (1) and (2) are the same as (4) and (5), respectively.

$$w^* = \arg\min_{w} \|y - Aw\|_1, \tag{4}$$

$$z^* = \underset{z}{\arg\min} \|z\|_1$$
 subject to  $x = Dz$ . (5)

### 3-D ball in different $\ell^r$ norms and the constraint Dz = x



3-D ball in the different  $\ell^r$  norms for r=2, 1.5, 1, 0.5

$$z^* = \underset{z}{\operatorname{arg\,min}} \|z\|_1$$
 subject to  $\underbrace{x}_{\text{given}} = D\underbrace{z}_{\text{many}}$  (5)

### The sparse representation problem

• We have introduced some ideas about the  $\ell^1$ - $\ell^0$  equivalence. In what follows, we don't consider the original SR problem. We consider the following  $\ell^1$  minimization problem instead:

**SR problem:** Given a signal vector  $x \in \mathbb{R}^m$  and a dictionary matrix  $D \in \mathbb{R}^{m \times n}$ , we seek a coefficient vector  $z^* \in \mathbb{R}^n$  such that

$$z^* = \arg\min_{z} \left(\frac{1}{2} \|x - Dz\|_2^2 + \lambda \|z\|_1\right), \quad \lambda > 0.$$
 (\*)

The existence (and uniqueness) of solution of the SR problem  $(\star)$  can be ensured because matrix  $D^{\top}D$  is symmetric (+ positive *definite*) and the second term  $\lambda \| \cdot \|_1$  is a *convex function*.

- Problem  $(\star)$  is also a regression analysis method in statistics and machine learning. It is the so-called *least absolute shrinkage and* selection operator (LASSO).
  - R. J. Tibshirani, The lasso problem and uniqueness, *Electronic Journal of Statistics*, 7 (2013), pp. 1456-1490 ⊕ A. Ali, 13 (2019), pp. 2307-2347.

### Alternating direction method of multipliers (ADMM)

We will use the "Alternating Direction Method of Multipliers" to solve the above  $\ell^1$ -norm SR problem.

 ADMM is an iterative scheme for solving the following equality constrained convex optimization problems:

$$\min_{z} f(z)$$
 subject to  $Az = b$ .

- ADMM consists of three steps:
  - 1. adding an auxiliary variable **y** and a dual variable (multipliers) **v** and then scaled as **u**
  - 2. separating the new cost function into a sum of f(z) and g(y)
  - 3. using an iterative method to solve the problem
- Then the optimization problem can be re-posed as

$$\min_{z, y} (f(z) + g(y))$$
 subject to  $Az + By = c$ .

## Derivation of the ADMM: augmented Lagrangian

First, we formulate the augmented Lagrangian

$$L_{\rho}(z,y,v) := f(z) + g(y) + \underbrace{v^{\top}}_{multipliers} (Az + By - c) + \underbrace{\frac{\rho}{2} \|Az + By - c\|_2^2}_{penalty\ term},$$

where  $\rho > 0$  is the penalty parameter. Then the iterative scheme of the *augmented Lagrangian method (ALM)* is given by

$$(z^{(i+1)}, y^{(i+1)}) = \underset{z}{\operatorname{arg min}} L_{\rho}(z, y, v^{(i)}),$$
  
 $v^{(i+1)} = v^{(i)} + \rho(Az^{(i+1)} + By^{(i+1)} - c).$ 

In ADMM, z and y are updated in an alternating or sequential fashion, which accounts for the term *alternating direction*.

$$egin{array}{lcl} oldsymbol{z}^{(i+1)} &=& rg \min_{oldsymbol{z}} L_{
ho}(oldsymbol{z}, oldsymbol{y}^{(i)}, oldsymbol{v}^{(i)}), \ oldsymbol{y}^{(i+1)} &=& rg \min_{oldsymbol{y}} L_{
ho}(oldsymbol{z}^{(i+1)}, oldsymbol{y}, oldsymbol{v}^{(i)}), \ oldsymbol{v}^{(i+1)} &=& oldsymbol{v}^{(i)} + 
ho(oldsymbol{A} oldsymbol{z}^{(i+1)} + oldsymbol{B} oldsymbol{y}^{(i+1)} - c). \end{array}$$

## Scaled form of the augmented Lagrangian

The ADMM can be written in a slightly different form, which is often more convenient, by combining the linear and quadratic terms in the augmented Lagrangian and scaling the dual variable (multipliers) v.

Define the residual r := Az + By - c. Then

$$v^{\top}(Az + By - c) + \frac{\rho}{2} ||Az + By - c||_{2}^{2}$$

$$= v^{\top}r + \frac{\rho}{2} ||r||_{2}^{2} = \frac{\rho}{2} ||r + \frac{1}{\rho}v||_{2}^{2} - \frac{1}{2\rho} ||v||_{2}^{2}.$$

Set  $\pmb{u} = \frac{1}{\rho} \pmb{v}$ . Then  $L_{\rho}(\pmb{z}, \pmb{y}, \pmb{v}) = L_{\rho}(\pmb{z}, \pmb{y}, \pmb{u})$ , and

$$L_{\rho}(z,y,u) = f(z) + g(y) + \frac{\rho}{2} ||Az + By - c + u||_{2}^{2} - \frac{\rho}{2} ||u||_{2}^{2}.$$

#### **ADMM: scaled form**

The ADMM in the scaled form is given by

$$\begin{split} & z^{(i+1)} = \underset{z}{\arg\min} \Big( f(z) + g(\boldsymbol{y}^{(i)}) + \frac{\rho}{2} \|Az + B\boldsymbol{y}^{(i)} - c + \boldsymbol{u}^{(i)}\|_2^2 - \frac{\rho}{2} \|\boldsymbol{u}^{(i)}\|_2^2 \Big), \\ & y^{(i+1)} = \underset{y}{\arg\min} \Big( f(\boldsymbol{z}^{(i+1)}) + g(\boldsymbol{y}) + \frac{\rho}{2} \|Az^{(i+1)} + B\boldsymbol{y} - c + \boldsymbol{u}^{(i)}\|_2^2 - \frac{\rho}{2} \|\boldsymbol{u}^{(i)}\|_2^2 \Big), \\ & \boldsymbol{u}^{(i+1)} = \boldsymbol{u}^{(i)} + Az^{(i+1)} + B\boldsymbol{y}^{(i+1)} - c, \end{split}$$

where  $\rho > 0$  is the *penalty parameter* which is related to the convergent rate of the iterations.

Note that the terms in blue can be omitted in practical computations!

#### References:

- S. Boyd, N. Parikh, E. Chu, B. Peleato, and J. Eckstein, Distributed optimization and statistical learning via the ADMM, Foundations and Trends in Machine Learning, 3 (2010), pp. 1-122.
- ADMM算法原理詳解: https://zhuanlan.zhihu.com/p/448289351

# **ADMM** for the $\ell^1$ -norm SR problem

• For the  $\ell^1$ -norm SR problem,

$$z^* = \arg\min_{z} \left( \frac{1}{2} \|x - Dz\|_2^2 + \lambda \|z\|_1 \right), \quad \lambda > 0, \quad (\star)$$

we set

$$f(z) := \frac{1}{2} \|x - Dz\|_{2}^{2}, \ g(y) := \lambda \|y\|_{1}, \ Az + By = c \Leftrightarrow z - y = 0.$$

ullet The ADMM for the  $\ell^1$ -norm SR problem is given by

$$z^{(i+1)} = \arg\min_{z} \left(\frac{1}{2} \|x - Dz\|_{2}^{2} + \frac{\rho}{2} \|z - y^{(i)} + u^{(i)}\|_{2}^{2}\right), \quad (6_{1})$$

$$y^{(i+1)} = \arg\min_{z} \left(\lambda \|y\|_{1} + \frac{\rho}{2} \|z^{(i+1)} - y + u^{(i)}\|_{2}^{2}\right), \quad (6_{2})$$

$$u^{(i+1)} = u^{(i)} + z^{(i+1)} - y^{(i+1)}, \quad (6_{3})$$

where  $\rho > 0$  is *penalty parameter* related to the convergent rate of the iterations.

# **Solving minimization problem** (6<sub>1</sub>)

Define

$$F_1(z) := \frac{1}{2} \|x - Dz\|_2^2 + \frac{\rho}{2} \|z - y^{(i)} + u^{(i)}\|_2^2.$$

Then  $F_1$  is a quadratic function in variables  $z_1, z_2, \dots, z_n$  and  $F_1(z) \ge 0 \ \forall \ z \in \mathbb{R}^n$ . To solve " $\min_z F_1(z)$ ", first we compute

$$\nabla F_1(z) = -\mathbf{D}^\top (\mathbf{x} - \mathbf{D}z) + \rho \mathbf{I}(\mathbf{z} - \mathbf{y}^{(i)} + \mathbf{u}^{(i)}) 
= (\mathbf{D}^\top \mathbf{D} + \rho \mathbf{I})\mathbf{z} - (\mathbf{D}^\top \mathbf{x} + \rho (\mathbf{y}^{(i)} - \mathbf{u}^{(i)})).$$

Letting  $\nabla F_1(z) = \mathbf{0}$ , we have

$$(\boldsymbol{D}^{\top}\boldsymbol{D} + \rho \boldsymbol{I})\boldsymbol{z} = \big(\boldsymbol{D}^{\top}\boldsymbol{x} + \rho(\boldsymbol{y}^{(i)} - \boldsymbol{u}^{(i)})\big).$$

Therefore, we obtain the solution

$$\boldsymbol{z}^{(i+1)} = (\boldsymbol{D}^{\top} \boldsymbol{D} + \rho \boldsymbol{I})^{-1} (\boldsymbol{D}^{\top} \boldsymbol{x} + \rho (\boldsymbol{y}^{(i)} - \boldsymbol{u}^{(i)})).$$

# **Solving minimization problem** (62)

Using the *soft-thresholding function*  $S_{\lambda/\rho}$ , the solution of problem (6<sub>2</sub>) has the closed form (see next few pages):

$$\mathbf{y}^{(i+1)} = \mathcal{S}_{\lambda/\rho}(\mathbf{z}^{(i+1)} + \mathbf{u}^{(i)}),$$

where

$$S_{\lambda/\rho}(v) = \operatorname{sign}(v) \odot \max(\mathbf{0}, |v| - \lambda/\rho),$$

and  $sign(\cdot)$ ,  $max(\cdot, \cdot)$ , and  $|\cdot|$  are all applied to the input vector v component-wisely, and  $\odot$  is the Hadamard product.

Finally, the iterative scheme can be posed as follows:

$$\mathbf{z}^{(i+1)} = (\mathbf{D}^{\top} \mathbf{D} + \rho \mathbf{I})^{-1} (\mathbf{D}^{\top} \mathbf{x} + \rho (\mathbf{y}^{(i)} - \mathbf{u}^{(i)})),$$
 (7<sub>1</sub>)

$$y^{(i+1)} = S_{\lambda/\rho}(z^{(i+1)} + u^{(i)}),$$
 (7<sub>2</sub>)

$$u^{(i+1)} = u^{(i)} + z^{(i+1)} - y^{(i+1)}.$$
 (7<sub>3</sub>)

# **Details of the solution of problem** (62)

Recall the problem  $(6_2)$ ,

$$y^{(i+1)} = \arg\min_{y} \left( \lambda \|y\|_1 + \frac{\rho}{2} \|z^{(i+1)} - y + u^{(i)}\|_2^2 \right).$$
 (62)

Let 
$$v:=z^{(i+1)}+u^{(i)}\in\mathbb{R}^n$$
. Then we have 
$$y^{(i+1)}=\arg\min_{\pmb{y}}\Bigl(\lambda\|\pmb{y}\|_1+rac{\rho}{2}\|\pmb{v}-\pmb{y}\|_2^2\Bigr).$$

Define a real-valued function  $F_2(y)$  as follows:

$$F_{2}(y) = \lambda \|y\|_{1} + \frac{\rho}{2} \|v - y\|_{2}^{2}$$

$$= \left(\lambda |y_{1}| + \frac{\rho}{2} (v_{1} - y_{1})^{2}\right) + \dots + \left(\lambda |y_{n}| + \frac{\rho}{2} (v_{n} - y_{n})^{2}\right)$$

$$:= f_{1}(y_{1}) + \dots + f_{n}(y_{n}),$$

where we define

$$f_j(y) := \lambda |y| + \frac{\rho}{2} (v_j - y)^2 \quad \forall j = 1, 2, \cdots, n.$$

# Analysis of functions $f_i$

For simplicity of the presentation, we consider the function

$$f(y) = \lambda |y| + \frac{\rho}{2} (v - y)^2.$$

Computing the derivative of f(y) for  $y \neq 0$ , we have

$$f'(y) = \left\{ \begin{array}{ll} \lambda - \rho(v-y) & \forall \ y > 0, \\ -\lambda - \rho(v-y) & \forall \ y < 0. \end{array} \right.$$

Let f'(y) = 0. Then we have

$$y = v - \frac{\lambda}{\rho}$$
 for  $y > 0$  and  $y = v + \frac{\lambda}{\rho}$  for  $y < 0$ .

Therefore, the all critical numbers of f are given by

$$c=v-rac{\lambda}{
ho} ext{ if } c>0, \quad c=v+rac{\lambda}{
ho} ext{ if } c<0, \quad c=0.$$

*In order to find the minimum of f, we consider the following three cases:* 

$$v > \frac{\lambda}{\rho}, \quad v < -\frac{\lambda}{\rho}, \quad -\frac{\lambda}{\rho} \le v \le \frac{\lambda}{\rho}.$$

Case 1: 
$$v > \frac{\lambda}{\rho}$$

In this case,  $c=v-\frac{\lambda}{\rho}>0$  is the only critical number and

$$\begin{split} f(c) &= f(v - \frac{\lambda}{\rho}) &= \lambda \left(v - \frac{\lambda}{\rho}\right) + \frac{\rho}{2} \left(v - \left(v - \frac{\lambda}{\rho}\right)\right)^2 \\ &= \frac{\rho}{2} \left(v^2 - \left(v - \frac{\lambda}{\rho}\right)^2\right) < \frac{\rho}{2} v^2 = f(0). \end{split}$$

For  $y \ge 0$ , since f is a quadratic polynomial in y with positive leading coefficient, we can conclude that  $f(c) \le f(y)$  for all  $y \ge 0$ .

For y < 0, f(y) is monotone decreasing since

$$f'(y) = \lambda \operatorname{sign}(y) - \rho(v - y) = -\lambda - \rho v + \rho y$$
  
$$< -\lambda - \lambda + \rho y = -2\lambda + \rho y < 0,$$

which implies f(y) > f(0) for all y < 0.

*Therefore, f has a minimum at*  $c = v - \frac{\lambda}{\rho} > 0$ .

Case 2: 
$$v<-\frac{\lambda}{\rho}$$

In this case,  $c = v + \frac{\lambda}{\rho} < 0$  is the only critical number and

$$f(c) = f(v + \frac{\lambda}{\rho}) = -\lambda(v + \frac{\lambda}{\rho}) + \frac{\rho}{2}\left(v - (v + \frac{\lambda}{\rho})\right)^{2}$$
$$= \frac{\rho}{2}\left(v^{2} - (v + \frac{\lambda}{\rho})^{2}\right) < \frac{\rho}{2}v^{2} = f(0).$$

For  $y \le 0$ , since f is a quadratic polynomial in y with positive leading coefficient, we can conclude that  $f(c) \le f(y)$  for all  $y \le 0$ .

For y > 0, f(y) is monotone increasing since

$$f'(y) = \lambda \operatorname{sign}(y) - \rho(v - y) = \lambda - \rho v + \rho y$$
  
>  $\lambda + \lambda + \rho y = 2\lambda + \rho y > 0$ ,

which implies f(y) > f(0) for all y > 0.

Therefore, f has a minimum at  $c = v + \frac{\lambda}{\rho} > 0$ .

Case 3: 
$$-\frac{\lambda}{\rho} \le v \le \frac{\lambda}{\rho}$$

In this case, we have no critical number except the non-differentiable point y = 0.

For y > 0, we have

$$f'(y) = \lambda - \rho(v - y) = \lambda - \rho v + \rho y$$
  
  $\geq \lambda - \lambda + \rho y = \rho y > 0.$ 

Thus, f(y) is monotone increasing and then f(y) > f(0) for all y > 0. For y < 0, we have

$$f'(y) = -\lambda - \rho(v - y) = -\lambda - \rho v + \rho y$$
  
 
$$\leq -\lambda + \lambda + \rho y = \rho y < 0.$$

Thus, f(y) is monotone decreasing and then f(y) > f(0) for all y < 0. Therefore, f has a minimum at 0.

# **Solution of problem** (62)

By the above discussions, we have

$$\underset{y}{\operatorname{arg\,min}} f(y) = \begin{cases} v + \frac{\lambda}{\rho}, & \text{if } v < -\frac{\lambda}{\rho}, & \text{(case 2)} \\ 0, & \text{if } |v| \leq \frac{\lambda}{\rho}, & \text{(case 3)} \\ v - \frac{\lambda}{\rho}, & \text{if } v > \frac{\lambda}{\rho}. & \text{(case 1)} \end{cases}$$

In other words, we have

$$\underset{y}{\arg\min}f(y) = \mathcal{S}_{\lambda/\rho}(v) = \operatorname{sign}(v)\max(0,|v| - \lambda/\rho).$$

Therefore,

$$y^{(i+1)} = \underset{y}{\operatorname{arg\,min}} F_2(y) = \mathcal{S}_{\lambda/\rho}(v) = \mathcal{S}_{\lambda/\rho}(z^{(i+1)} + u^{(i)}).$$

where the soft-thresholding,

$$S_{\lambda/\rho}(v) := \operatorname{sign}(v) \odot \max(\mathbf{0}, |v| - \lambda/\rho),$$

and  $sign(\cdot)$ ,  $max(\cdot, \cdot)$ , and  $|\cdot|$  are all applied to the input vector v component-wisely, and  $\odot$  is the Hadamard product.

## Application to signal denoising

- First, we construct a random dictionary matrix  $D \in \mathbb{R}^{512 \times 2048}$  and a random sparse vector  $z \in \mathbb{R}^{2048}$  with  $||z||_0 = 32$ . We then have the true signal x := Dz.
- Define the noise signal  $x_n := x + n$ , where  $n \in \mathbb{R}^{512}$  is a random white Gaussian noise with noised powers P = 0.5, 1, 5 (噪聲功率). We consider  $\lambda = 5, 10, 20, 30$  for the minimization problem.
- Peak signal-to-noise ratio (PSNR, 峰值訊噪比): We define the mean squared error (MSE) and then the PSNR as follows:

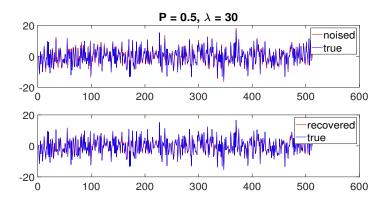
$$MSE := \frac{1}{512} \sum_{i=1}^{512} \left( \text{true}(i) - \text{approx}(i) \right)^2,$$

$$PSNR := 10 \times \log_{10} \left( \frac{\text{max}^2}{MSE} \right),$$

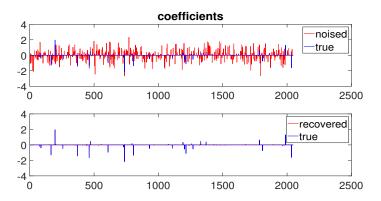
where "max" is the maximum amplitude of the true signal x.

• Source of matlab code: http://brendt.wohlberg.net/software/SPORCO/

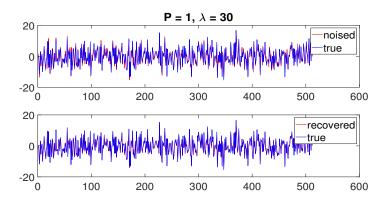
#### Numerical results for P = 0.5 and $\lambda = 30$



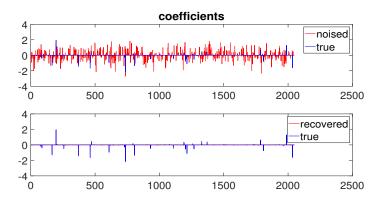
### **Coefficients for** P = 0.5 **and** $\lambda = 30$



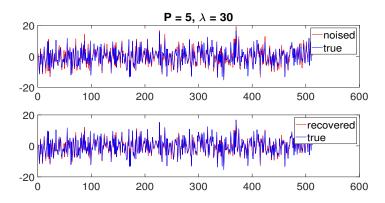
### Numerical results for P = 1 and $\lambda = 30$



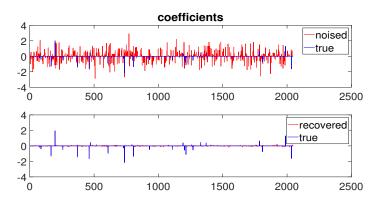
### Coefficients for P = 1 and $\lambda = 30$



### Numerical result for P = 5 and $\lambda = 30$



### Coefficients for P = 5 and $\lambda = 30$



#### **PSNR** values and iteration numbers

*In general, the higher the value of PSNR the better the quality of the recovered signals.* 

#### **PSNR** values

P	0.5	0.5	1	1	5	5
λ	noised	rcvered	noised	rcvered	noised	rcvered
5	29.51	30.36	29.71	30.41	25.57	26.11
10	29.51	31.16	29.71	31.10	25.57	26.63
20	29.51	32.55	29.71	32.23	25.57	27.62
30	29.51	33.45	29.71	32.77	25.57	28.50

#### Iteration numbers of ADMM

$\lambda P$	0.5	1	5
5	550	664	569
10	301	303	320
20	172	169	186
30	129	130	154

## Sparse dictionary learning problem

In the SR problem, the solution of interest  $z^*$  is the coefficient vector of a linear combination of over-complete basis elements (columns) from a given dictionary  $\mathbf{D}$  under some sparsity constraint. Therefore, it is typically accompanied by a dictionary learning mechanism.

We are going to study a more general problem. The dictionary D is unknown and needed to be sought together with the sparse solution.

**SDL problem:** Let  $\{x_i\}_{i=1}^N \subset \mathbb{R}^m$  be a given dataset of signals. We seek a dictionary matrix  $\mathbf{D} = [d_1, d_2, \cdots, d_n] \in \mathbb{R}^{m \times n}$  together with the sparse coefficient vectors  $\{z_i\}_{i=1}^N \subset \mathbb{R}^n$  that solve the minimization problem:

$$\begin{split} \min_{D,\{z_i\}} \left( \frac{1}{2} \sum_{i=1}^N \|x_i - Dz_i\|_2^2 + \lambda \sum_{i=1}^N \|z_i\|_1 \right) \\ \text{subject to } \|d_k\|_2 \leq 1, \ \forall \ 1 \leq k \leq n, \end{split}$$

where  $\lambda > 0$  is a penalty parameter.

## Problem formulation in a more compact form

To simplify the formulation of the SDL problem, we define

$$X = [x_1, x_2, \cdots, x_N] \in \mathbb{R}^{m \times N},$$
  
 $Z = [z_1, z_2, \cdots, z_N] \in \mathbb{R}^{n \times N}.$ 

Then the SDL problem can be posed as follows: Given a training data matrix X, find a dictionary matrix D and a coefficient matrix Z such that

$$\min_{D, Z} \left( \frac{1}{2} \| X - DZ \|_F^2 + \lambda \| Z \|_{1,1} \right) \tag{$\star \star$}$$
 subject to  $\| d_k \|_2 \le 1, \ \forall \ 1 \le k \le n,$ 

where  $\|\cdot\|_F$  denotes the Frobenius norm and  $\|\mathbf{Z}\|_{1,1}$  is the  $L_{1,1}$ -norm which is defined as

$$\|\mathbf{Z}\|_{1,1} := \sum_{i=1}^{N} \|z_i\|_1.$$

### An iterative approach for solving the SDL problem

In the SDL problem  $(\star\star)$ , we have two unknown matrices D and Z. We will use a simple iterative approach together with the ADMM to solve  $(\star\star)$ , though it is more complicated.

Given an initial guess  $D_{(0)}$ , for  $j=0,1,\cdots$ , we solve the following two sub-problems alternatingly:

$$Z_{(j)} = \arg\min_{\mathbf{Z}} \left( \frac{1}{2} \| \mathbf{X} - \mathbf{D}_{(j)} \mathbf{Z} \|_{F}^{2} + \lambda \| \mathbf{Z} \|_{1,1} \right),$$
(8)  

$$D_{(j+1)} = \arg\min_{\mathbf{D}} \left( \frac{1}{2} \| \mathbf{X} - \mathbf{D} \mathbf{Z}_{(j)} \|_{F}^{2} + \lambda \| \mathbf{Z}_{(j)} \|_{1,1} \right)$$
subject to  $\| \mathbf{d}_{k} \|_{2} \le 1, \ \forall \ 1 \le k \le n.$ (9)

We iterate (8) and (9) until convergence is achieved. As we have introduced previously, problems (8) and (9) will be solved by ADMM.

## **ADMM** for solving problem (8)

• Adding an auxiliary variable Y and a dual variable U, we define

$$f(\mathbf{Z}) := \frac{1}{2} \|\mathbf{X} - \mathbf{D}_{(j)} \mathbf{Z}\|_F^2, \quad g(\mathbf{Y}) := \lambda \|\mathbf{Y}\|_{1,1}, \quad \mathbf{Z} = \mathbf{Y}.$$

• Then the ADMM for solving (8) is given by

$$Z^{(i+1)} = \underset{\mathbf{Z}}{\operatorname{arg\,min}} \left( \frac{1}{2} \| \mathbf{X} - \mathbf{D}_{(j)} \mathbf{Z} \|_F^2 + \frac{\rho}{2} \| \mathbf{Z} - \mathbf{Y}^{(i)} + \mathbf{U}^{(i)} \|_F^2 \right), \quad (8_1)$$

$$Y^{(i+1)} = \underset{\mathbf{Y}}{\operatorname{arg\,min}} \left( \lambda \| \mathbf{Y} \|_{1,1} + \frac{\rho}{2} \| \mathbf{Z}^{(i+1)} - \mathbf{Y} + \mathbf{U}^{(i)} \|_F^2 \right), \quad (8_2)$$

$$\mathbf{U}^{(i+1)} = \mathbf{U}^{(i)} + \mathbf{Z}^{(i+1)} - \mathbf{Y}^{(i+1)}. \quad (8_3).$$

• Similar to the SR problem, we will use the same methods to solve the sub-problems  $(8_1)$  and  $(8_2)$ .

# **Solving minimization problem** $(8_1)$

Define

$$F_1(\mathbf{Z}) := \frac{1}{2} \| \mathbf{X} - \mathbf{D}_{(j)} \mathbf{Z} \|_F^2 + \frac{\rho}{2} \| \mathbf{Z} - \mathbf{Y}^{(i)} + \mathbf{U}^{(i)} \|_F^2.$$

To solve "min  $F_1(\mathbf{Z})$ ", first we compute

$$\nabla F_1(\mathbf{Z}) = -\mathbf{D}_{(j)}^{\top} (\mathbf{X} - \mathbf{D}_{(j)} \mathbf{Z}) + \rho \mathbf{I} (\mathbf{Z} - \mathbf{Y}^{(i)} + \mathbf{U}^{(i)})$$
$$= (\mathbf{D}_{(j)}^{\top} \mathbf{D}_{(j)} + \rho \mathbf{I}) \mathbf{Z} - (\mathbf{D}_{(j)}^{\top} \mathbf{X} + \rho (\mathbf{Y}^{(i)} - \mathbf{U}^{(i)})).$$

Letting  $\nabla F_1(\mathbf{Z}) = 0$ , we have

$$(\boldsymbol{D}_{(j)}^{\top}\boldsymbol{D}_{(j)}+\rho\boldsymbol{I})\boldsymbol{Z}=\big(\boldsymbol{D}_{(j)}^{\top}\boldsymbol{X}+\rho(\boldsymbol{Y}^{(i)}-\boldsymbol{U}^{(i)})\big).$$

Therefore, we obtain the solution

$$\mathbf{Z}^{(i+1)} = (\mathbf{D}_{(i)}^{\top} \mathbf{D}_{(i)} + \rho \mathbf{I})^{-1} (\mathbf{D}_{(i)}^{\top} \mathbf{X} + \rho (\mathbf{Y}^{(i)} - \mathbf{U}^{(i)})).$$

# **Solving minimization problem** (8<sub>2</sub>)

Using the component-wise soft-thresholding function, the solution of problem  $(8_2)$  has the closed form:

$$\mathbf{Y}^{(i+1)} = \mathcal{S}_{\lambda/\rho}(\mathbf{Z}^{(i+1)} + \mathbf{U}^{(i)}),$$

where

$$S_{\lambda/\rho}(V) = \operatorname{sign}(V) \odot \max(\mathbf{0}, |V| - \lambda/\rho),$$

with sign(V) and |V| are element-wisely applied to the matrix V and  $\odot$  is the Hadamard product.

Therefore, the iterative scheme can be posed as follows:

$$\mathbf{Z}^{(i+1)} = (\mathbf{D}_{(j)}^{\top} \mathbf{D}_{(j)} + \rho \mathbf{I})^{-1} (\mathbf{D}_{(j)}^{\top} \mathbf{X} + \rho (\mathbf{Y}^{(i)} - \mathbf{U}^{(i)})), \qquad (10_1)$$

$$\mathbf{Y}^{(i+1)} = \mathcal{S}_{\lambda/\rho} (\mathbf{Z}^{(i+1)} + \mathbf{U}^{(i)}), \qquad (10_2)$$

$$\mathbf{U}^{(i+1)} = \mathbf{U}^{(i)} + \mathbf{Z}^{(i+1)} - \mathbf{Y}^{(i+1)}. \qquad (10_3)$$

## **Solving minimization problem (9)**

Recall that

$$D_{(j+1)} = \arg\min_{D} \left( \frac{1}{2} \| X - DZ_{(j)} \|_{F}^{2} + \lambda \| Z_{(j)} \|_{1,1} \right)$$
subject to  $\| d_{k} \|_{2} \le 1$ ,  $\forall 1 \le k \le n$ . (9)

Since the term  $\lambda \|\mathbf{Z}_{(j)}\|_{1,1}$  is a fixed number when  $\mathbf{Z}_{(j)}$  is given, problem (9) can be replaced by

$$D_{(j+1)} = \underset{D}{\operatorname{arg min}} \frac{1}{2} ||X - DZ_{(j)}||_F^2$$
  
subject to  $||d_k||_2 \le 1$ ,  $\forall 1 \le k \le n$ . (9')

Next, we introduce an *auxiliary variable* G and a *dual variable* H in ADMM for solving (9').

### ADMM for solving problem (9')

Define

$$g(G) := \{[d_1, d_2, \cdots, d_n] : ||d_k||_2 \le 1, \forall 1 \le k \le n\},\ G := D.$$

The ADMM for solving problem (9') is given by

$$\begin{split} & \boldsymbol{D}^{(i+1)} &= & \arg\min_{\boldsymbol{D}} \left( \frac{1}{2} \| \boldsymbol{X} - \boldsymbol{D} \boldsymbol{Z}_{(j)} \|_F^2 + \frac{\rho}{2} \| \boldsymbol{D} - \boldsymbol{G}^{(i)} + \boldsymbol{H}^{(i)} \|_F^2 \right), \quad (9_1) \\ & \boldsymbol{G}^{(i+1)} &= & \operatorname{proj}_{\mathcal{G}(G)} \{ \boldsymbol{D}^{(i+1)} \}, \quad (9_2) \\ & \boldsymbol{H}^{(i+1)} &= & \boldsymbol{H}^{(i)} + \boldsymbol{D}^{(i+1)} - \boldsymbol{G}^{(i+1)}. \quad (9_3) \end{split}$$

For solving problem  $(9_1)$ , we define

$$F_2(\mathbf{D}) := \frac{1}{2} \| \mathbf{X} - \mathbf{D} \mathbf{Z}_{(j)} \|_F^2 + \frac{\rho}{2} \| \mathbf{D} - \mathbf{G}^{(i)} + \mathbf{H}^{(i)} \|_F^2.$$

# **Solving minimization problem** (9<sub>1</sub>)

Computing  $\nabla F_2(\mathbf{D})$ , we have

$$\nabla F_2(\mathbf{D}) = (\mathbf{X} - \mathbf{D}\mathbf{Z}_{(j)})(-\mathbf{Z}_{(j)}^{\top}) + \rho I_m(\mathbf{D} - \mathbf{G}^{(i)} + \mathbf{H}^{(i)})$$
  
=  $\mathbf{D}(\rho I_n + \mathbf{Z}_{(j)}\mathbf{Z}_{(j)}^{\top}) + \mathbf{X}\mathbf{Z}_{(j)}^{\top} - \rho(\mathbf{G}^{(i)} - \mathbf{H}^{(i)}).$ 

Letting  $\nabla F_2(\mathbf{D}) = \mathbf{0}$ , we have

$$D(\mathbf{Z}_{(j)}\mathbf{Z}_{(j)}^{\top} + \rho \mathbf{I}_n) = \mathbf{X}\mathbf{Z}_{(j)}^{\top} - \rho(\mathbf{G}^{(i)} - \mathbf{H}^{(i)}).$$

Therefore, we obtain the solution

$$\boldsymbol{D}^{(i+1)} = (\boldsymbol{X} \boldsymbol{Z}_{(j)}^{\top} - \rho (\boldsymbol{G}^{(i)} - \boldsymbol{H}^{(i)})) (\boldsymbol{Z}_{(j)} \boldsymbol{Z}_{(j)}^{\top} + \rho \boldsymbol{I}_n)^{-1}.$$

Finally, the ADMM for problem (9') is given by

$$\begin{split} & \boldsymbol{D}^{(i+1)} &= (\boldsymbol{X} \boldsymbol{Z}_{(j)}^{\top} - \rho (\boldsymbol{G}^{(i)} - \boldsymbol{H}^{(i)})) (\boldsymbol{Z}_{(j)} \boldsymbol{Z}_{(j)}^{\top} + \rho \boldsymbol{I}_n)^{-1}, \quad (11_1) \\ & \boldsymbol{G}^{(i+1)} &= \operatorname{proj}_{\boldsymbol{g}(\boldsymbol{G})} \{\boldsymbol{D}^{(i+1)}\}, \quad (11_2) \\ & \boldsymbol{H}^{(i+1)} &= \boldsymbol{H}^{(i)} + \boldsymbol{D}^{(i+1)} - \boldsymbol{G}^{(i+1)}. \quad (11_3) \end{split}$$

#### Convergence and stopping criterion

- In [Boyd *et al.* 2010], there are more details about convergence results of the ADMM.
- In the iterative scheme  $(10_1)$ ,  $(10_2)$ ,  $(10_3)$ , we define

$$R_z = Y^{(i+1)} - Y^{(i)}, \quad S_z = U^{(i+1)} - U^{(i)}.$$

If  $R_z$  and  $S_z$  less than the tolerances  $\varepsilon_{R_z}$  and  $\varepsilon_{S_z}$ , then we say that the iteration of coefficients  $\mathbf{Z}^{(i+1)}$  converges.

In the iterative scheme  $(11_1)$ ,  $(11_2)$ ,  $(11_3)$ , we define

$$R_d = G^{(i+1)} - G^{(i)}, \quad S_d = H^{(i+1)} - H^{(i)}.$$

If  $R_d$  and  $S_d$  less than the tolerances  $\varepsilon_{R_d}$  and  $\varepsilon_{S_d}$ , then we say that the iteration of dictionary  $D^{(i+1)}$  converges.

#### References and source codes

# $\ell^1$ - $\ell^0$ equivalence problem:

- [1] D. L. Donoho, For most large underdetermined systems of linear equations the minimal  $\ell_1$ -norm solution is also the sparsest solution, *Communications on Pure and Applied Mathematics*, 59 (2006), pp. 797-829.
- [2] Y. Sharon, J. Wright, and Y. Ma, Computation and relaxation of conditions for equivalence between  $\ell^1$  and  $\ell^0$  minimization, *UIUC Technical Report UILU-ENG-07-2008*, 2007.

#### ADMM:

[1] S. Boyd, N. Parikh, E. Chu, B. Peleato, and J. Eckstein, Distributed optimization and statistical learning via the ADMM, *Foundations and Trends in Machine Learning*, 3 (2010), pp. 1-122.

#### Sparse dictionary learning:

https://en.wikipedia.org/wiki/Sparse\_dictionary\_learning

Matlab codes: http://brendt.wohlberg.net/software/SPORCO/

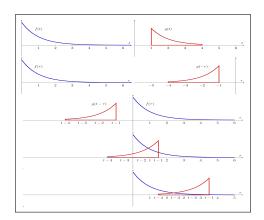
#### Part II

Convolutional Sparse Representation and Dictionary Learning

#### **Convolution of two functions**

Let f and g be two integrable functions with compact supports in  $\mathbb{R}$ . Then the convolution of f and g is defined as a function in variable t,

$$(f * g)(t) := \int_{-\infty}^{\infty} f(\tau)g(t - \tau) d\tau, \quad t \in \mathbb{R}.$$



#### **Convolution of two vectors**

**Definition:** Let  $u = [u_1, \dots, u_n]^{\top} \in \mathbb{R}^n$  and  $v = [v_1, \dots, v_m]^{\top} \in \mathbb{R}^m$ . The convolution of u and v, denoted by u \* v, is defined as follows:

$$\boldsymbol{u} * \boldsymbol{v} := \begin{bmatrix} u_1 v_1 \\ u_1 v_2 + u_2 v_1 \\ u_1 v_3 + u_2 v_2 + u_3 v_1 \\ \vdots \\ u_{n-2} v_m + u_{n-1} v_{m-1} + u_n v_{m-2} \\ u_{n-1} v_m + u_n v_{m-1} \\ u_n v_m \end{bmatrix} \in \mathbb{R}^{m+n-1}.$$

More specifically, for  $i = 1, 2, \dots, (m + n - 1)$ , the i-th component of u \* v is given by

$$(\boldsymbol{u} * \boldsymbol{v})_i = \sum_{j=\max(1,i-m+1)}^{\min(i,n)} u_j v_{i-j+1}.$$

**Remark:** Convolutional operator \* is commutative, i.e., u \* v = v \* u.

## Convolutional sparse representation (CSR) problem

**CSR problem:** Given a signal  $\mathbf{z} \in \mathbb{R}^m$  and a dictionary  $\mathbf{D} = [\mathbf{d}_1, \cdots, \mathbf{d}_n] \in \mathbb{R}^{\ell \times n}$ , we seek a sparse matrix  $\mathbf{Z} = [\mathbf{z}_1, \cdots, \mathbf{z}_n] \in \mathbb{R}^{k \times n}$ ,  $m = \ell + k - 1$ , which solves the following minimization problem:

$$\min_{\mathbf{Z}} \left( \frac{1}{2} \| \mathbf{x} - \sum_{j=1}^{n} d_j * z_j \|_2^2 + \lambda \sum_{j=1}^{n} \| z_j \|_1 \right),$$

where  $\lambda > 0$  is a penalty parameter.

#### **Remarks:**

• In SR, we use Dz to recover the signal x,

$$x \approx Dz = d_1z_1 + d_2z_2 + \cdots + d_nz_n = \sum_{j=1}^n d_jz_j.$$

In CSR, we use  $\sum_{j=1}^{n} d_j * z_j$  instead,

$$x \approx d_1 * z_1 + d_2 * z_2 + \cdots + d_n * z_n = \sum_{j=1}^n d_j * z_j.$$

• Convolution is a way to regulate  $d_j * z_j$  such that  $x \approx \sum_{j=1}^n d_j * z_j$ . It is more flexible than  $x \approx \sum_{j=1}^n d_j z_j$ , but indeed more expensive!

## **Toeplitz matrix**

We define an  $(m + n - 1) \times m$  matrix *U* in terms of  $u_i$ , which is called a *Toeplitz matrix*, as follows:

$$\boldsymbol{U} := \begin{bmatrix} u_1 & 0 & \cdots & 0 & 0 \\ u_2 & u_1 & \ddots & 0 & 0 \\ \vdots & u_2 & \ddots & 0 & 0 \\ \vdots & \vdots & \ddots & u_1 & 0 \\ u_{n-1} & \vdots & \vdots & u_2 & u_1 \\ u_n & u_{n-1} & \vdots & \vdots & u_2 \\ 0 & u_n & \ddots & \vdots & \vdots \\ 0 & 0 & \ddots & u_{n-1} & \vdots \\ \vdots & \vdots & \ddots & u_n & u_{n-1} \\ 0 & 0 & \cdots & 0 & u_n \end{bmatrix}_{(m+n-1)\times m}^{\top} \in \mathbb{R}^n$$
Then one can check that  $\boldsymbol{u} * \boldsymbol{v} = \boldsymbol{U} \boldsymbol{v}$ , where  $\boldsymbol{u} = [u_1, \cdots, u_n]^{\top} \in \mathbb{R}^n$  and  $\boldsymbol{v} = [v_1, \cdots, v_m]^{\top} \in \mathbb{R}^m$ .

and  $v = [v_1, \cdots, v_m]^\top \in \mathbb{R}^m$ .

## **CSR** problem using Toeplitz matrices

With the help of Toeplitz matrix, we can rewrite the CSR problem as

$$\min_{\widetilde{z}} \left( \frac{1}{2} \| x - \widetilde{D} \widetilde{z} \|_{2}^{2} + \lambda \| \widetilde{z} \|_{1} \right), \tag{12}$$

with

$$\widetilde{z} = [z_1^\top, z_2^\top, \cdots, z_n^\top]_{nk \times 1}^\top$$
 and  $\widetilde{D} = [D_1, D_2, \cdots, D_n]_{(\ell+k-1) \times nk}$ 

where  $\mathbf{D}_j$  is a Toeplitz  $(\ell + k - 1) \times k$  matrix associated with the column vector  $\mathbf{d}_i \in \mathbb{R}^{\ell}$ , and  $\ell + k - 1 = m$ .

#### Remarks:

- We can use the same way for SR problem to solve the CSR problem (12). We can employ the ADMM, but it is too expensive since the matrix size of  $\widetilde{D}$  is too large.
- The discrete Fourier transform  $\mathcal{F}: \mathbb{C}^N \to \mathbb{C}^N$  can help us to address this computational issue.

#### Discrete Fourier transform (DFT) and its inverse (IDFT)

•  $\widehat{\mathbf{x}} = \mathcal{F}(\mathbf{x})$ : The DFT  $\mathcal{F}: \mathbb{C}^N \to \mathbb{C}^N$  transforms a finite vector  $\mathbf{x} = [x_1, x_2, \cdots, x_N]^{\top}$  into another vector  $\hat{\mathbf{x}} = [\hat{x}_1, \hat{x}_2, \cdots, \hat{x}_N]^{\top}$ , which is defined by

$$\widehat{x}_k = \sum_{n=1}^N x_n e^{-\frac{i2\pi}{N}(k-1)(n-1)}.$$

Then DFT is an invertible linear transformation.

•  $x = \mathcal{F}^{-1}(\widehat{x})$ : The inverse discrete Fourier transform (IDFT)  $\mathcal{F}^{-1}: \mathbb{C}^N \to \mathbb{C}^N, \widehat{x} \mapsto x$ , is given by

$$x_n = \frac{1}{N} \sum_{k=1}^{N} \widehat{x}_k e^{\frac{i2\pi}{N}(k-1)(n-1)}.$$

• Euler's formula:  $e^{i\theta} = \cos \theta + i \sin \theta$ ,  $\forall \theta \in \mathbb{R}$ .

https://en.wikipedia.org/wiki/Discrete\_Fourier\_transform

## Hadamard product

• Let  $u \in \mathbb{R}^n$  and  $v \in \mathbb{R}^m$ . Then  $u * v \in \mathbb{R}^{m+n-1}$  and

$$\mathcal{F}(u*v) = \mathcal{F}(u') \odot \mathcal{F}(v'),$$

where  $\mathcal{F}$  denotes the DFT, u' and v' are respectively the zero padding of u and v with the same size of u \* v, i.e.,

$$u' = [u^{\top}, 0, \dots, 0]^{\top}, \quad v' = [v^{\top}, 0, \dots, 0]^{\top} \in \mathbb{R}^{m+n-1},$$

and  $\odot$  is the Hadamard product.

• The Hadamard product  $\odot$  of two vectors is a component-wise product. Let  $\mathbf{u} = [u_1, u_2, \cdots, u_n]^\top, \mathbf{v} = [v_1, v_2, \cdots, v_n]^\top \in \mathbb{R}^n$ ,

$$\boldsymbol{u} \odot \boldsymbol{v} := [u_1 v_1, u_2 v_2, \cdots, u_n v_n]^{\top}.$$

We can define a diagonal matrix U such that  $u \odot v = Uv$ , where

$$\boldsymbol{U} := \begin{bmatrix} u_1 & 0 & \cdots & 0 \\ 0 & u_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & u_n \end{bmatrix}.$$

# Recalling the CSR problem

**CSR problem:** Given  $x \in \mathbb{R}^m$  and  $D = [d_1, \dots, d_n] \in \mathbb{R}^{\ell \times n}$ , we seek  $Z = [z_1, \dots, z_n] \in \mathbb{R}^{k \times n}$  with  $m = \ell + k - 1$  solving

$$\min_{\mathbf{Z}} \left( \frac{1}{2} \| \mathbf{x} - \sum_{j=1}^{n} d_j * \mathbf{z}_j \|_2^2 + \lambda \sum_{j=1}^{n} \| \mathbf{z}_j \|_1 \right).$$

To solve the above minimization problem, we first use the ADMM algorithm to split it into three subproblems:

$$\begin{split} & \boldsymbol{Z}^{(i+1)} &= & \arg\min_{\boldsymbol{Z}} \left( \frac{1}{2} \| \boldsymbol{x} - \sum_{j=1}^{n} \boldsymbol{d}_{j} * \boldsymbol{z}_{j} \|_{2}^{2} + \frac{\rho}{2} \sum_{j=1}^{n} \| \boldsymbol{z}_{j} - \boldsymbol{y}_{j}^{(i)} + \boldsymbol{u}_{j}^{(i)} \|_{2}^{2} \right), \\ & \boldsymbol{Y}^{(i+1)} &= & \arg\min_{\boldsymbol{Y}} \left( \lambda \sum_{j=1}^{n} \| \boldsymbol{y}_{j} \|_{1} + \frac{\rho}{2} \sum_{j=1}^{n} \| \boldsymbol{z}_{j}^{(i+1)} - \boldsymbol{y}_{j} + \boldsymbol{u}_{j}^{(i)} \|_{2}^{2}, \right) \\ & \boldsymbol{U}^{(i+1)} &= & \boldsymbol{U}^{(i)} + \boldsymbol{Z}^{(i+1)} - \boldsymbol{Y}^{(i+1)}. \end{split}$$

# Using discrete Fourier transform for Z

We will use the discrete Fourier transform and Hadamard product to solve the subproblem of **Z**. We can rewrite these subproblems as

$$\begin{split} \widehat{\mathbf{Z}}^{(i+1)} &= & \arg\min_{\widehat{\mathbf{Z}}} \left(\frac{1}{2} \| \widehat{\mathbf{x}} - \sum_{j=1}^n \widehat{d_j'} \odot \widehat{\mathbf{z}_j'} \|_2^2 + \frac{\rho}{2} \sum_{j=1}^n \| \widehat{\mathbf{z}_j'} - \widehat{\mathbf{y}_j'}^{(i)} + \widehat{\mathbf{u}_j'}^{(i)} \|_2^2 \right), \\ \mathbf{Y}^{(i+1)} &= & \arg\min_{\mathbf{Y}} \left(\lambda \sum_{j=1}^n \| \mathbf{y}_j \|_1 + \frac{\rho}{2} \sum_{j=1}^n \| \mathcal{F}^{-1}(\widehat{\mathbf{z}_j'}^{(i+1)}) - \mathbf{y}_j + \mathbf{u}_j^{(i)} \|_2^2 \right), \\ \mathbf{U}^{(i+1)} &= & \mathbf{U}^{(i)} + \mathcal{F}^{-1}(\widehat{\mathbf{Z}}^{(i+1)}) - \mathbf{Y}^{(i+1)}, \\ \text{where} & \\ & \mathcal{F}^{-1}(\widehat{\mathbf{Z}}) = [\mathcal{F}^{-1}(\widehat{\mathbf{z}_1'}), \mathcal{F}^{-1}(\widehat{\mathbf{z}_2'}), \cdots, \mathcal{F}^{-1}(\widehat{\mathbf{z}_n'})]. \end{split}$$

#### Why can we use the discrete Fourier transform?

Note that the discrete Fourier transform  $\mathcal{F}$  is linear. Thus, we have

$$\frac{1}{2} \| \mathbf{x} - \sum_{j=1}^{n} d_{j} * \mathbf{z}_{j} \|_{2}^{2} = \frac{1}{2m} \| \mathcal{F}(\mathbf{x} - \sum_{j=1}^{n} d_{j} * \mathbf{z}_{j}) \|_{2}^{2} \quad \text{(Plancherel theorem)}$$

$$= \frac{1}{2m} \| \mathcal{F}(\mathbf{x}) - \mathcal{F}(\sum_{j=1}^{n} d_{j} * \mathbf{z}_{j}) \|_{2}^{2}$$

$$= \frac{1}{2m} \| \widehat{\mathbf{x}} - \sum_{j=1}^{n} \mathcal{F}(d_{j} * \mathbf{z}_{j}) \|_{2}^{2} = \frac{1}{2m} \| \widehat{\mathbf{x}} - \sum_{j=1}^{n} \widehat{d}'_{j} \odot \widehat{\mathbf{z}}'_{j} \|_{2}^{2}.$$

Similarly, the second term of subproblem **Z** can be rewritten as

$$\frac{\rho}{2} \sum_{j=1}^{n} \| z_j - y_j^{(i)} + u_j^{(i)} \|_2^2 = \frac{\rho}{2} \sum_{j=1}^{n} \| z_j' - y_j'^{(i)} + u_j'^{(i)} \|_2^2 
= \frac{\rho}{2m} \sum_{i=1}^{n} \| \widehat{z_j'} - \widehat{y_j'}^{(i)} + \widehat{u_j'}^{(i)} \|_2^2.$$

**Note:**  $x \in \mathbb{R}^m$ ,  $d_j \in \mathbb{R}^\ell$ ,  $z_j \in \mathbb{R}^k$ ,  $d_j * z_j \in \mathbb{R}^{\ell+k-1} = \mathbb{R}^m$ ,  $d_j'$ ,  $z_j' \in \mathbb{R}^m$ .

# The subproblem of $\hat{Z}$

We first define

$$\widehat{m{D}_j} = \mathrm{diag}(\widehat{m{d}_j'}). \quad (m \times m \ disgonal \ matrix)$$

Then the subproblem in the Fourier domain can be posed as:

$$\widehat{\boldsymbol{z}}^{(i+1)} = \arg\min_{\widehat{\boldsymbol{z}}} \left( \frac{1}{2} \| \widehat{\boldsymbol{x}} - \widehat{\boldsymbol{D}} \widehat{\boldsymbol{z}} \|_2^2 + \frac{\rho}{2} \| \widehat{\boldsymbol{z}} - \widehat{\boldsymbol{y}}^{(i)} + \widehat{\boldsymbol{u}}^{(i)} \|_2^2 \right),$$

where

$$\widehat{D} = [\widehat{D}_1, \widehat{D}_2, \cdots, \widehat{D}_n]_{m \times mn}, \quad \widehat{z} = [\widehat{z_1^{\prime}}^{\top}, \widehat{z_2^{\prime}}^{\top}, \cdots, \widehat{z_n^{\prime}}^{\top}]_{mn \times 1}^{\top},$$

and

$$\widehat{y} = [\widehat{y_1'}^\top, \widehat{y_2'}^\top, \cdots, \widehat{y_n'}^\top]_{mn \times 1}^\top, \quad \widehat{u} = [\widehat{u_1'}^\top, \widehat{u_2'}^\top, \cdots, \widehat{u_n'}^\top]_{mn \times 1}^\top.$$

*Note that we drop the scalar factor 1/m in the subproblem.* 

### Rewriting the subproblems of ADMM

Using the above definitions, we can rewrite the subproblems of ADMM as:

$$\widehat{z}^{(i+1)} = \min_{\widehat{z}} \left( \frac{1}{2} \| \widehat{x} - \widehat{D} \widehat{z} \|_{2}^{2} + \frac{\rho}{2} \| \widehat{z} - \widehat{y}^{(i)} + \widehat{u}^{(i)} \|_{2}^{2} \right), \quad (13_{1})$$

$$y^{(i+1)} = \min_{y} \left( \lambda \| y \|_{1} + \frac{\rho}{2} \sum_{j=1}^{n} \| \mathcal{F}^{-1}(\widehat{z}^{(i+1)}) - y + u^{(i)} \|_{2}^{2} \right), \quad (13_{2})$$

$$u^{(i+1)} = u^{(i)} + \mathcal{F}^{-1}(\widehat{z}^{(i+1)}) - y^{(i+1)}, \quad (13_{3})$$

where

$$y = [{y_1'}^{\top}, {y_2'}^{\top}, \cdots, {y_n'}^{\top}]_{mn \times 1}^{\top}, \quad u = [{u_1'}^{\top}, {u_2'}^{\top}, \cdots, {u_n'}^{\top}]_{mn \times 1}^{\top}.$$

# **Solving minimization problem** (13<sub>1</sub>)

First we define

$$F(\widehat{z}) = \frac{1}{2} \|\widehat{x} - \widehat{D}\widehat{z}\|_{2}^{2} + \frac{\rho}{2} \|\widehat{z} - \widehat{y}^{(i)} + \widehat{u}^{(i)}\|_{2}^{2}.$$

To solve  $\min_{\widehat{z}} F(\widehat{z})$ , we compute

$$\nabla F(\widehat{z}) = -\widehat{D}^{\top}(\widehat{x} - \widehat{D}\widehat{z}) + \rho I(\widehat{z} - \widehat{y}^{(i)} + \widehat{u}^{(i)})$$
$$= (\widehat{D}^{\top}\widehat{D} + \rho I)\widehat{z} - (\widehat{D}^{\top}\widehat{x} + \rho(\widehat{y}^{(i)} - \widehat{u}^{(i)})).$$

Letting  $\nabla F(\widehat{z}) = \mathbf{0}$ , we have

$$(\widehat{\boldsymbol{D}}^{\top}\widehat{\boldsymbol{D}} + \rho \boldsymbol{I})_{mn \times mn}\widehat{\boldsymbol{z}} = \widehat{\boldsymbol{D}}^{\top}\widehat{\boldsymbol{x}} + \rho(\widehat{\boldsymbol{y}}^{(i)} - \widehat{\boldsymbol{u}}^{(i)}).$$

Therefore, we obtain the solution

$$\widehat{\boldsymbol{z}}^{(i+1)} = (\widehat{\boldsymbol{D}}^{\top} \widehat{\boldsymbol{D}} + \rho \boldsymbol{I})^{-1} \big(\widehat{\boldsymbol{D}}^{\top} \widehat{\boldsymbol{x}} + \rho (\widehat{\boldsymbol{y}}^{(i)} - \widehat{\boldsymbol{u}}^{(i)})\big).$$

# **Solving minimization problem** (13<sub>2</sub>)

The way to solve minimization problem  $(13_2)$  is similar to that for solving problem  $(6_2)$ .

Finally, we obtain the ADMM iterative scheme as follows:

$$\widehat{\mathbf{z}}^{(i+1)} = (\widehat{\mathbf{D}}^{\top} \widehat{\mathbf{D}} + \rho \mathbf{I})^{-1} (\widehat{\mathbf{D}}^{\top} \widehat{\mathbf{x}} + \rho (\widehat{\mathbf{y}}^{(i)} - \widehat{\mathbf{u}}^{(i)})), \quad (14_1)$$

$$\mathbf{y}^{(i+1)} = \mathcal{S}_{\lambda/\rho} (\mathcal{F}^{-1} (\widehat{\mathbf{z}}^{(i+1)}) + \mathbf{u}^{(i)}), \quad (14_2)$$

$$\mathbf{u}^{(i+1)} = \mathbf{u}^{(i)} + \mathcal{F}^{-1} (\widehat{\mathbf{z}}^{(i+1)}) - \mathbf{y}^{(i+1)}. \quad (14_3)$$

Next, we will introduce the *Sherman-Morrison formula* which can be applied to solve  $\hat{z}^{(i+1)}$  in a more efficient way.

## The Sherman-Morrison formula: a special case

Let A and B be two  $n \times n$  matrices. In general, (A + B) is not invertible, even though A is invertible. However, if A is invertible and B has some certain structure, then  $(A + B)^{-1}$  exists.

**A special case of the Sherman-Morrison formula:** *Let* I *be the*  $n \times n$  *identity matrix and* u, v *be two given vectors in*  $\mathbb{C}^n$ . *If*  $1 + v^{\top}u \neq 0$ , *then*  $I + uv^{\top}$  *is invertible and* 

$$(I + uv^{\top})^{-1} = I - \frac{uv^{\top}}{1 + v^{\top}u}.$$

Proof: We check that

$$(I + uv^{\top}) \left( I - \frac{uv^{\top}}{1 + v^{\top}u} \right) = I - \frac{uv^{\top}}{1 + v^{\top}u} + uv^{\top} - \frac{uv^{\top}uv^{\top}}{1 + v^{\top}u}$$

$$= I + \frac{-uv^{\top} + uv^{\top} + v^{\top}uuv^{\top}}{1 + v^{\top}u} - \frac{uv^{\top}uv^{\top}}{1 + v^{\top}u}$$

$$= I + \frac{v^{\top}uuv^{\top}}{1 + v^{\top}u} - \frac{uv^{\top}uv^{\top}}{1 + v^{\top}u} = I. \quad (v^{\top}u : scalar) \quad \Box$$

#### How to derive the inverse?

Given  $b \in \mathbb{C}^n$ , we consider the linear system  $(I + uv^\top)x = b$ . Assume that  $I + uv^\top$  is invertible. Then the unique solution x exists. Let  $k = v^\top x \in \mathbb{C}$ . Then  $x + ku = b \Rightarrow v^\top x + kv^\top u = v^\top b$   $\Rightarrow k + k(v^\top u) = v^\top b$ , which implies

$$k = \frac{\boldsymbol{v}^{\top} \boldsymbol{b}}{1 + \boldsymbol{v}^{\top} \boldsymbol{u}}, \quad \text{if } 1 + \boldsymbol{v}^{\top} \boldsymbol{u} \neq 0.$$

Therefore, we know that

$$x = b - ku = b - \frac{v^{\top}b}{1 + v^{\top}u}u = b - \frac{uv^{\top}}{1 + v^{\top}u}b = \underbrace{\left(I - \frac{uv^{\top}}{1 + v^{\top}u}\right)}_{\text{inverse of }I + uv^{\top}}b.$$

**The Sherman-Morrison formula:** Suppose that  $A \in \mathbb{C}^{n \times n}$  is an invertible matrix and  $u, v \in \mathbb{C}^n$ . Then  $A + uv^{\top}$  is invertible if and only if  $1 + v^{\top}A^{-1}u \neq 0$ . In this case, we have

$$(A + uv^{\top})^{-1} = A^{-1} - \frac{A^{-1}uv^{\top}A^{-1}}{1 + v^{\top}A^{-1}u}.$$

# How to compute $(14_1)$ ?

Recall that

$$(\widehat{\boldsymbol{D}}^{\top}\widehat{\boldsymbol{D}} + \rho \boldsymbol{I})\widehat{\boldsymbol{z}}^{(i+1)} = (\widehat{\boldsymbol{D}}^{\top}\widehat{\boldsymbol{x}} + \rho(\widehat{\boldsymbol{y}}^{(i)} - \widehat{\boldsymbol{u}}^{(i)})), \quad (14_1)$$

where matrix  $\hat{D}$  has the following structure:

$$\widehat{D} = [\widehat{D}_{1}, \widehat{D}_{2}, \cdots, \widehat{D}_{n}]_{m \times mn}$$

$$= \begin{bmatrix} \widehat{d'_{1,1}} & 0 & \cdots & 0 & \widehat{d'_{2,1}} & 0 & \cdots & 0 & \cdots \\ 0 & \widehat{d'_{1,2}} & \ddots & \vdots & 0 & \widehat{d'_{2,2}} & \ddots & \vdots & \cdots \\ \vdots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & 0 & \cdots \\ 0 & \cdots & 0 & \widehat{d'_{1,m}} & 0 & \cdots & 0 & \widehat{d'_{2,m}} & \cdots \end{bmatrix},$$

where

$$\widehat{D_j} = \mathrm{diag}(\widehat{d'_j}) \quad (m \times m \ diagonal \ matrix).$$

# The structure of matrix $\widehat{\boldsymbol{D}}^{\top}\widehat{\boldsymbol{D}} + \rho \boldsymbol{I}$

Note that

$$\begin{split} \widehat{D}^{\top}\widehat{D} + \rho I &= \begin{bmatrix} \widehat{D}_{1}^{\top} \\ \widehat{D}_{2}^{\top} \\ \vdots \\ \widehat{D}_{n}^{\top} \end{bmatrix} [\widehat{D}_{1}, \widehat{D}_{2}, \cdots, \widehat{D}_{n}] + \rho I \\ &= \begin{bmatrix} \widehat{D}_{1}^{\top}\widehat{D}_{1} + \rho I_{m} & \widehat{D}_{1}^{\top}\widehat{D}_{2} & \cdots & \widehat{D}_{1}^{\top}\widehat{D}_{n} \\ \widehat{D}_{2}^{\top}\widehat{D}_{1} & \widehat{D}_{2}^{\top}\widehat{D}_{2} + \rho I_{m} & \cdots & \widehat{D}_{2}^{\top}\widehat{D}_{n} \\ \vdots & \vdots & \ddots & \vdots \\ \widehat{D}_{n}^{\top}\widehat{D}_{1} & \widehat{D}_{n}^{\top}\widehat{D}_{2} & \cdots & \widehat{D}_{n}^{\top}\widehat{D}_{n} + \rho I_{m} \end{bmatrix}_{mn \times mn} \end{split}$$

By re-ordering the equations, the  $mn \times mn$  system  $(\widehat{D}^{\top}\widehat{D} + \rho I)\widehat{z}^{(i+1)} = R$ can be replaced by m independent linear systems of size  $n \times n$ , each of which consists of a rank one component plus a diagonal component, then solved by the Sherman-Morrison formula, see [Wohlberg 2016, Appendix A].

# Convolutional sparse dictionary learning problem

We now consider the convolutional sparse dictionary learning problem, where the dictionary  $\boldsymbol{D}$  is unknown and needed to be sought together with the convolutional sparse solution.

**Convolutional SDL problem:** Let  $\{x_i\}_{i=1}^N \subset \mathbb{R}^m$  be a given dataset of signals. We seek a dictionary matrix  $\mathbf{D} = [\mathbf{d}_1, \mathbf{d}_2, \cdots, \mathbf{d}_M] \in \mathbb{R}^{\ell \times M}$  and the coefficient matrices  $\{\mathbf{Z}_i\}_{i=1}^N \subset \mathbb{R}^{k \times M}$  with  $\mathbf{Z}_i = [\mathbf{z}_{1,i}, \mathbf{z}_{2,i}, \cdots, \mathbf{z}_{M,i}]$  and  $m = \ell + k - 1$  such that  $\mathbf{D}$  and  $\{\mathbf{Z}_i\}_{i=1}^N$  solve the following minimization problem:

$$\begin{split} \min_{\{d_j\}_{j=1}^M, \{z_{j,i}\}_{j=1,i=1}^{M,N}} \left(\frac{1}{2} \sum_{i=1}^N \|x_i - \sum_{j=1}^M d_j * z_{j,i}\|_2^2 + \lambda \sum_{i=1}^N \sum_{j=1}^M \|z_{j,i}\|_1 \right) \\ \text{subject to } \|d_j\|_2 \leq 1 \quad \forall \, j = 1, 2, \cdots, M, \end{split}$$

where  $\lambda > 0$  is a given penalty parameter.

#### **Toeplitz matrix**

• Define  $\widetilde{D} = \begin{bmatrix} D_1 & D_2 & \cdots & D_M \end{bmatrix}$  with each  $D_j$  is the *Toeplitz matrix* defined with respect to  $d_j$  as before.

Define 
$$z_i = \begin{bmatrix} z_{1,i}^\top & z_{2,i}^\top & \cdots & z_{M,i}^\top \end{bmatrix}^\top$$
 for  $i = 1, 2, \cdots, N$ , where each  $z_i$  is the coefficient vector with respect to the data  $x_i$ . Define  $\mathbf{Z} = \begin{bmatrix} z_1 & z_2 & \cdots & z_N \end{bmatrix}$  and  $\mathbf{X} = \begin{bmatrix} x_1 & x_2 & \cdots & x_N \end{bmatrix}$ .

• The convolutional SDL problem can be simplified as

$$\min_{\widetilde{D}, Z} \left( \frac{1}{2} \| X - \widetilde{D} Z \|_F^2 + \lambda \| Z \|_{1,1} \right)$$
subject to  $\| d_j \|_2 \le 1 \quad \forall j = 1, 2, \dots, M$ ,

where  $\|\mathbf{Z}\|_{1,1}$  is defined as before.

#### How to solve the convolutional SDL problem?

Though we can still use the ADMM iterative scheme to solve

$$\min_{\widetilde{D}, Z} \left( \frac{1}{2} \| X - \widetilde{D}Z \|_F^2 + \lambda \| Z \|_{1,1} \right)$$
subject to  $\| d_j \|_2 \le 1 \quad \forall j = 1, 2, \cdots, M$ ,

the sizes of the involved matrices are too large. Thus, we will use the DFT and the Sherman-Morrison formula to deal with this problem. The steps are similar to the CSR problem, but more complicated.

Recall the convolutional SDL problem:

$$\begin{split} \min_{\{d_j\}_{j=1}^M, \{z_{j,i}\}_{j=1,i=1}^{M,N}} \left(\frac{1}{2} \sum_{i=1}^N \|x_i - \sum_{j=1}^M d_j * z_{j,i}\|_2^2 + \lambda \sum_{i=1}^N \sum_{j=1}^M \|z_{j,i}\|_1 \right) \\ \text{subject to } \|d_j\|_2 \leq 1 \quad \forall \ j = 1, 2, \cdots, M. \end{split}$$

For solving this problem, we split it into two parts.

## Step 1: Solving the coefficient Z

For solving the convolutional SDL problem, we first give an initial dictionary  $D = [d_1, d_2, \cdots, d_M]$  to solve coefficient Z and then further use ADMM algorithm to split this problem into three subproblems:

$$\widehat{Z}^{(i+1)} = \arg\min_{\widehat{Z}} \left( \frac{1}{2} \sum_{i=1}^{N} \| \widehat{x}_i - \sum_{j=1}^{M} \widehat{d}'_j \odot \widehat{z'_{j,i}} \|_2^2 + \frac{\rho}{2} \sum_{i=1}^{N} \sum_{j=1}^{M} \| \widehat{z'_{j,i}} - \widehat{y'_{j,i}}^{(i)} + \widehat{u'_{j,i}}^{(i)} \|_2^2 \right)$$

$$\mathbf{Y}^{(i+1)} = \arg\min_{\mathbf{Y}} \Big( \lambda \sum_{i=1}^{N} \sum_{j=1}^{M} \|\mathbf{y}_{j,i}\|_{1} + \frac{\rho}{2} \sum_{i=1}^{N} \sum_{j=1}^{M} \|\mathcal{F}^{-1}(\widehat{\mathbf{z}'_{j,i}}) - \mathbf{y}_{j,i} + \mathbf{u}'_{j,i}\|_{2}^{2} \Big),$$

$$\mathbf{U}^{(i+1)} = \mathbf{U}^{(i)} + \mathcal{F}^{-1}(\widehat{\mathbf{Z}}^{(i+1)}) - \mathbf{Y}^{(i+1)}$$

with

$$\mathcal{F}^{-1}(\widehat{\mathbf{Z}}) = \begin{bmatrix} \mathcal{F}^{-1}(\widehat{z_{1,1}'}) & \cdots & \mathcal{F}^{-1}(\widehat{z_{1,M}'}) \\ \vdots & \ddots & \vdots \\ \mathcal{F}^{-1}(\widehat{z_{N,1}'}) & \cdots & \mathcal{F}^{-1}(\widehat{z_{N,M}'}) \end{bmatrix}.$$

### Rewritten in a compact form

For convenience, we can rewrite these subproblems as follows:

$$\begin{split} \widehat{\pmb{Z}}^{(i+1)} &= & \arg\min_{\widehat{\pmb{Z}}} \left(\frac{1}{2} \| \widehat{\pmb{X}} - \widehat{\pmb{D}} \widehat{\pmb{Z}} \|_F^2 + \frac{\rho}{2} \| \widehat{\pmb{Z}} - \widehat{\pmb{Y}}^{(i)} + \widehat{\pmb{U}}^{(i)} \|_F^2 \right), \\ \pmb{Y}^{(i+1)} &= & \arg\min_{\pmb{Y}} \left( \lambda \| \pmb{Y} \|_{1,1} + \frac{\rho}{2} \| \mathcal{F}^{-1} (\widehat{\pmb{Z}}^{(i+1)}) - \widehat{\pmb{Y}} + \widehat{\pmb{U}}^{(i)} \|_F^2 \right), \\ \pmb{U}^{(i+1)} &= & \pmb{U}^{(i)} + \mathcal{F}^{-1} (\widehat{\pmb{Z}}^{(i+1)}) - \pmb{Y}^{(i+1)}, \end{split}$$

with

$$\widehat{X} = [\widehat{x_1}, \widehat{x_2}, \cdots, \widehat{x_N}], \quad \widehat{D} = [\widehat{D_1}, \widehat{D_2}, \cdots, \widehat{D_M}],$$

and

$$\mathbf{Y} = \begin{bmatrix} y_{1,1}' & \cdots & y_{1,M}' \\ \vdots & \ddots & \vdots \\ y_{N,1}' & \cdots & y_{N,M}' \end{bmatrix}, \quad \mathbf{U} = \begin{bmatrix} u_{1,1}' & \cdots & u_{1,M}' \\ \vdots & \ddots & \vdots \\ u_{N,1}' & \cdots & u_{N,M}' \end{bmatrix}.$$

Using the similar ways as that for solving CSR problem, we can solve the above subproblems.

### **Step 2: Solving the dictionary** *D*

Recall the convolutional SDL problem:

$$\min_{\substack{\{d_j\}_{j=1}^M,\ \{z_{j,i}\}_{j=1,i=1}^M \\ }} \left(\frac{1}{2} \sum_{i=1}^N \|x_i - \sum_{j=1}^M d_j * z_{j,i}\|_2^2 + \lambda \sum_{i=1}^N \sum_{j=1}^M \|z_{j,i}\|_1\right)$$
subject to  $\|d_j\|_2 \le 1 \quad \forall j = 1, 2, \cdots, M$ .

When the coefficient Z is obtained, the blue term is a given number. Solving the dictionary D is equivalent to solve

$$\min_{\{d_j\}_{j=1}^M} \frac{1}{2} \sum_{i=1}^N \|x_i - \sum_{j=1}^M d_j * z_{j,i}\|_2^2 \text{ subject to } \|d_j\|_2 \le 1, \quad \forall j = 1, 2, \cdots, M.$$

#### Using ADMM algorithm to solve Step 2

We use the ADMM algorithm to solve the above problem:

$$\begin{split} & \boldsymbol{D}^{(i+1)} &= & \arg\min_{\boldsymbol{D}} \Big( \frac{1}{2} \sum_{i=1}^{N} \| \boldsymbol{x}_i - \sum_{j=1}^{M} \boldsymbol{d}_j * \boldsymbol{z}_j \|_2^2 + \frac{\rho}{2} \sum_{j=1}^{M} \| \boldsymbol{d}_j - \boldsymbol{g}_j^{(i)} + \boldsymbol{h}_j^{(i)} \|_2^2 \Big), \\ & \boldsymbol{G}^{(i+1)} &= & \operatorname{proj}_{\boldsymbol{g}_{(\boldsymbol{G})}} \{ \boldsymbol{D}^{(i+1)} \}, \\ & \boldsymbol{H}^{(i+1)} &= & \boldsymbol{H}^{(i)} + \boldsymbol{D}^{(i+1)} - \boldsymbol{G}^{(i+1)}, \end{split}$$

and then use the Fourier transform and similar ways as before,

$$\begin{split} \widehat{\boldsymbol{D}}^{(i+1)} &= \arg\min_{\widehat{\boldsymbol{D}}} \Big(\frac{1}{2} \sum_{i=1}^{N} \|\widehat{\boldsymbol{x}}_{i} - \sum_{j=1}^{M} \widehat{\boldsymbol{d}}_{j}^{\prime} \odot \widehat{\boldsymbol{z}}_{j}^{\prime} \|_{2}^{2} + \frac{\rho}{2} \sum_{j=1}^{M} \|\widehat{\boldsymbol{d}}_{j}^{\prime} - \widehat{\boldsymbol{g}}_{j}^{\prime(i)} + \widehat{\boldsymbol{h}}_{j}^{\prime(i)} \|_{2}^{2} \Big), \\ G^{(i+1)} &= \operatorname{proj}_{\boldsymbol{g}_{(G)}} \{ \mathcal{F}^{-1}(\widehat{\boldsymbol{D}}^{(i+1)}) \}, \\ \boldsymbol{H}^{(i+1)} &= \boldsymbol{H}^{(i)} + \mathcal{F}^{-1}(\widehat{\boldsymbol{D}}^{(i+1)}) - \boldsymbol{G}^{(i+1)}, \\ \text{where} \\ \mathcal{F}^{-1}(\widehat{\boldsymbol{D}}) &= [\mathcal{F}^{-1}(\widehat{\boldsymbol{d}}_{1}^{\prime}), \mathcal{F}^{-1}(\widehat{\boldsymbol{d}}_{2}^{\prime}), \cdots, \mathcal{F}^{-1}(\widehat{\boldsymbol{d}}_{M}^{\prime}) ]. \end{split}$$

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